

Answers to End of Chapter Reviews and Exercises

for Assembly Language for x86 Processors, 8th Edition

by Kip R. Irvine

Chapters 1 to 13

Revision date: 3/10/2019

Chapter 1

1.7.1 Short Answer Questions

1. Most significant bit (the highest numbered bit).
2. (a) 53 (b) 150 (c) 204
3. (a) 110001010 (b) 110010110 (c) 100100001
4. 00000110
5. (a) 8 (b) 32 (c) 64 (d) 128
6. (a) 12 (b) 16 (c) 16
7. (a) 35DA (b) CEA3 (c) FEDB
8. (a) 0000 0001 0010 0110 1111 1001 1101 0100
(b) 0110 1010 1100 1101 1111 1010 1001 0101
(c) 1111 0110 1001 1011 1101 1100 0010 1010
9. (a) 58 (b) 447 (c) 16534
10. (a) 98 (b) 1203 (c) 671
11. (a) FFE8 (b) FEB5
12. (a) FFEB (b) FFD3
13. (a) 27641 (b) ?16093
14. (a) 19666 (b) ?32208
15. (a) -75 (b) +42 (c) -16
16. (a) -128 (b) -52 (c) -73
17. (a) 11111011 (b) 11010110 (c) 11110000
18. (a) 10111000 (b) 10011110 (c) 11100110
19. (a) AB2 (b) 1106
20. (a) B82 (b) 1316
21. 42h and 66d
22. 47h and 71d

23. $2^{29} - 1$, or 6.8056473384187692692674921486353 X 10^{38}

24. $2^{86} - 1$, or 77371252455336267181195263

25. Truth table:

A	B	$A \vee B$	$\neg(A \vee B)$
F	F	F	T
F	T	T	F
T	F	T	F
T	T	T	F

26. Truth table: (last column is the same as #25)

A	B	$\neg A$	$\neg B$	$\neg A \wedge \neg B$
F	F	T	T	T
F	T	T	F	F
T	F	F	T	F
T	T	F	F	F

27. It requires 2^4 (16) rows.

28. 2 bits, producing the following values: 00, 01, 10, 11

1.7.2 Algorithm Workbench

1. Code example (C++)

```
int toInt32(string s) {
    int num = 0;
    for(int i = 0; s[i] >= '0' && s[i] <= '1'; i++) {
        num = num * 2 + s[i] - '0';
    }
    return num;
}
```

2. Code example (C++)

```
int hexStrToInt32(string s) {
    int num = 0;
    for(int i = 0; ; i++) {
        if( s[i] >= '0' && s[i] <= '9' )
            num = num * 16 + s[i] - '0';
        else if( s[i] >= 'A' && s[i] <= 'F' )
            num = num * 16 + (s[i] - 'A' + 10);
        else
            break;
    }
    return num;
}
```

3. Code example (C++)

```
string intToBinStr( int n ) {
    vector<int> stack;

    do {
        int quotient = n / 2;
        int remainder = n % 2;
        stack.push_back(remainder);
        n = quotient;
    } while( n > 0 );

    string s;
    while( stack.size() > 0 ) {
        s += (stack.back() + '0');
        stack.pop_back();
    }
    return s;
}
```

4. Code example (C++)

```
string intToHexStr( int n ) {
    vector<int> stack;

    do {
        int quotient = n / 16;
        int remainder = n % 16;
        stack.push_back(remainder);
        n = quotient;
    } while( n > 0 );

    string s;
    while( stack.size() > 0 ) {
        int d = stack.back();
        if( d >= 0 && d <= 9 )
            s += (stack.back() + '0');
        else // probably a hex digit
            s += (stack.back() - 10 + 'A');
        stack.pop_back();
    }
    return s;
}
```

5. Code example (C++)

```
string addDigitStrings( string s1, string s2, int base ) {
    string sumStr;
    int carry = 0;

    for(int i = s1.size() - 1; i >= 0; i--) {
        int dval = (s1[i] - '0') + (s2[i] - '0') + carry;
        carry = 0;
        if( dval > (base - 1) ) {
            carry = 1;
            dval = dval % base;
        }
        sumStr.insert(sumStr.begin(), (dval + '0'));
    }
    if( carry == 1 )
        sumStr.insert( sumStr.begin(), 1 + '0');
    return sumStr;
}
```

Chapter 2

2.8.1 Short Answer Questions

1. EBP
2. Choose 4 from: Carry, Zero, Sign, Direction, Aux Carry, Overflow, Parity.
3. Carry flag
4. Overflow flag
5. True
6. Sign flag
7. Floating-point unit
8. 80 bits
9. True
10. False
11. True
12. False
13. True
14. False
14. False
15. False
16. True
17. False
18. True
19. False
20. False
21. True
22. True
23. False
24. False
25. Hardware, BIOS, and OS
26. It gives them more precise control of hardware, and execution is faster.

Chapter 3

3.9.1 Short Answer Questions

1. ADD, SUB, MOV
2. A calling convention determines how parameters are passed to subroutines, and how the stack is restored after the subroutine call.
3. By subtracting a value from the stack pointer register.
4. Assembler means the program that translates your source code. A more correct term is "assembly language".
5. Little endian places the least significant bit in position 0, on the right side of the number. Big endian does the opposite.
6. An integer literal, such as 35, has no direct meaning to someone reading the program's source code. Instead, a symbolic constant such as MIN_CAPACITY can be assigned an integer value, and is self-documenting.
7. A source file is given as input to the assembler. A listing file has additional text that will not assemble. It is a file that is created by the assembler.
8. Data labels exist in the data segment as variable offsets. Code labels are in the code segment, and are offsets for transfer of control instructions.
9. True
10. False (this notation is used in C, but not in assembly language).
11. False
12. True
13. Label, mnemonic, operand(s), comment
14. True
15. True
16. Code example:

```
COMMENT !
    this is the first comment line
    this is the second comment line
!
```
17. You do not use numeric addresses (offsets) for variables because the addresses would change if new variables were inserted before the existing ones.

3.9.2 Algorithm Workbench

1. Code example:

```
one = 25
two = 11001b
three = 31o
four = 19h
```
2. Yes, it can have multiple code and data segments.

3. Storing the value 01020304h

```
myVal LABEL DWORD
BYTE 04h,03h,02h,01h
```

4. Yes, you can. The assembler does not check the number's sign.

5. Code example. The two instructions have different opcodes.

```
add eax,5
add edx,5
```

6. Little endian order: ABh, 89h, 67h, 45h

7. myArray DWORD 120 DUP(?)

8. firstFive BYTE "ABCDE"

9. smallVal SDWORD 80000000h ; -32,768

10. wArray WORD 1000h, 2000h, 3000h

11. favColor BYTE "blue",0

12. dArray DWORD 50 DUP(?)

13. msg BYTE 500 DUP("TEST")

14. bArray BYTE 20 DUP(0)

Chapter 4

4.9.1 Short Answer

1. a. `edx = FFFF8002h` b. `edx = 00004321h`
2. `eax = 10020000h`
3. `eax = 3002FFFFh`
4. `eax = 10020001h`
5. Parity Even (1)
6. `eax = FFFFFFFFh`, `SF = 1` (the result is negative)
7. $-1 + 130 = 129$, which is outside the range of a signed positive byte. Therefore, the Overflow flag is set.
8. `rax = 0000000044445555h`
9. `rax = FFFFFFFF84326732h`
10. `eax = 00035678h`
11. `eax = 12375678h`
12. No.
13. Yes.
14. Yes, for example:

```
mov al,-128
neg al      ; OF = 1
```
15. No.
16. (a) not valid, (b) valid, (c) not valid, (d) not valid, (e) not valid, (f) not valid, (g) valid, (h) not valid
17. (a) `FCh` (b) `01h`
18. (a) `1000h`, (b) `3000h`, (c) `FFF0h`, (d) `4000h`
19. (a) `00000001h`, (b) `00001000h`, (c) `00000002h`, (d) `FFFFFFFCh`

4.9.2 Algorithm Workbench

1. Code example:

```
mov ax,word ptr three
mov bx,word ptr three+2
mov three,bx
mov word ptr three+2,ax
```
2. Code example: convert `al,bl,cl,dl` to `bl,cl,dl,al`

```
xchg al,bl
xchg al,cl
xchg al,dl
```

3. Code example:

```
mov al,01110101
add al,0           ; PF = 0 (odd)
```

4. Code example:

```
mov al,-127
add al,-1         ; OF = 1
```

5. Code example (set Zero and Carry)

```
mov al,0FFh
add al,1
```

6. Code example:

```
mov al,3
sub al,4
```

7. Code example: $AX = (val2 + BX) - val4$

```
mov ax,val2
add ax,bx
sub ax,val4
```

8. Code example:

```
mov al,80h
add al,80h
```

9. Code example:

```
mov ax,val2
neg ax
add ax,bx
sub ax,val4
```

10. Setting the Carry and Overflow flags at the same time:

```
mov al,80h
add al,80h
```

11. Setting the Zero flag after INC and DEC to indicate unsigned overflow:

```
mov al,0FFh
inc al
jz overflow_occurred
mov bl,1
dec bl
jz overflow_occurred
```

12. Data directive:

```
.data
ALIGN 2
myBytes BYTE 10h, 20h, 30h, 40h
etc.
```

13. (a) 1 (b) 4 (c) 4 (d) 2 (e) 4 (f) 8 (g) 5

14. Code:

```
mov dx, WORD PTR myBytes
```

15. Code:

```
mov al, BYTE PTR myWords+1
```


16. Code:

```
mov eax, DWORD PTR myBytes
```

17. Data directive:

```
myWordsD LABEL DWORD  
myWords WORD 3 DUP(?),2000h  
.code  
mov eax,myWordsD
```

18. Data directive:

```
myBytesW LABEL WORD  
myBytes BYTE 10h,20h,30h,40h  
.code  
mov ax,myBytesW
```

Chapter 5

5.7.1 Short Answer

1. pusha
2. pushf
3. popf
4. Because you might not want to push all the general-purpose registers when EAX is being used to pass a return value back to the subroutine's caller.

5. Code example (32-bit mode):

```
sub esp,4
mov [esp],eax
```

6. True
7. False
8. True
9. False
10. Yes, the pointer is in the ESI register
11. True
12. False
13. False
14. The following statements would have to be modified:

```
mov [esi],eax      becomes --> mov [esi],ax
add esi,4          becomes --> add esi,2
```

15. EAX = 5
16. (d)
17. (c)
18. (c)
19. (a)
20. The array will contain 10, 20, 30, 40

5.7.2 Algorithm Workbench

1. Exchange registers using push and pop:

```
push ebx
push eax
pop ebx
pop eax
```

2. Modify a subroutine's return address:

```
pop    eax                ; get the return address
add    eax,3              ; add 3
push   eax                ; put it back on the stack
ret
```

3. Create and assign local variables:

```
sub    esp,8              ; space for two variables
mov    [esp],1000h
mov    [esp+4],2000h
```

4. Copy an array element backwards

```
mov    edi,esi
dec    edi
mov    edx,array[esi*4]
mov    array[edi*4],edx
```

5. Display a subroutine's return address

```
mov    eax,[esp]
call   WriteHex
```

Chapter 6

6.10.1 Short Answer

1. BX = 006Bh
2. BX = 092h
3. BX = 064BBh
4. BX = A857h
5. EBX = BFAFF69Fh
6. RBX = 0000000050509B64h
7. AL = 2Dh, 48h, 6Fh, A3h
8. AL = 85h, 34h, BFh, AEh
9. a. CF= 0 ZF= 0 SF=0
b. CF= 0 ZF= 0 SF=0
c. CF= 1 ZF= 0 SF=1
10. JECX
11. JA and JNB jump to the destination if ZF = 0 and CF = 0.
12. EDX = 1
13. EDX = 1
14. EDX = 0
15. True
16. True
17. 0FFFFFFFFFFFFFFF80h
18. 0FFFFFFFF808080h
19. 0000000080808080h

6.10.2 Algorithm Workbench

1. `and al,0Fh`
2. Calculate parity of a doubleword:

```
.data
memVal DWORD ?
.code
    mov al,BYTE PTR memVal
    xor al,BYTE PTR memVal+1
    xor al,BYTE PTR memVal+2
    xor al,BYTE PTR memVal+3
```

3. Generate a bit string in EAX that represents members in SetX that are not members of SetY:

```
.data
SetX DWORD ?
SetY DWORD ?
.code
    mov  eax,SetX
    xor  eax,SetY      ; remove all SetY from SetX
```

4. Jump to label L1 when the unsigned integer in DX is less than or equal to the integer in CX:

```
    cmp  dx,cx
    jbe  L1
L1:
```

5. Write instructions that jump to label L2 when the signed integer in AX is greater than the

integer in CX:

```
    cmp  ax,cx
    jg   L2
L2:
```

6. First clear bits 0 and 1 in AL. Then, if the destination operand is equal to zero, the code should jump to label L3. Otherwise, it should jump to label L4:

```
    and  al,00000011b
    jz   L3
    jmp  L4
    mov  edx,0
L3:
L4:
```

7. Code example:

```
    cmp  val1,cx
    jna  L1
    cmp  cx,dx
    jna  L1
    mov  X,1
    jmp  next
L1: mov  X,2
next:
```

8. Code example:

```
    cmp  bx,cx
    ja   L1
    cmp  bx,val1
    ja   L1
    mov  X,2
    jmp  next
L1: mov  X,1
next:
```

9. Code example:

```
    cmp bx,cx          ; bx > cx?
    jna L1            ; no: try condition after OR
    cmp bx,dx         ; yes: is bx > dx?
    jna L1            ; no: try condition after OR
    jmp L2            ; yes: set X to 1

;-----OR(dx > ax) -----
L1:  cmp dx,ax        ; dx > ax?
    jna L3            ; no: set X to 2

L2:  mov X,1          ; yes:set X to 1
    jmp next         ; and quit

L3:  mov X,2          ; set X to 2
next:
```

10. Code example:

```
Exercise10Test proc
; use these registers to hold the logical variables:
    mov  eax,4        ; A
    mov  ebx,5        ; B
    mov  edx,10       ; N
    call Exercise10Test
    ret
Exercise10Test endp

Exercise10 proc
whileloop:
    cmp  edx,0
    jle  endwhile
    cmp  edx,3        ; if N != 3
    je   elselabel
    ; check N < eax OR N > ebx
    cmp  edx,eax     ; N < A?
    jl   orlabel     ; if true, jump
    cmp  edx,ebx     ; or N > B?
    jg   orlabel     ; if true, jump
    jmp  elselabel
orlabel:
    sub  edx,2
    jmp  whileloop
elselabel:
    sub  edx,1
    jmp  whileloop
endwhile:
    ret
Exercise10 endp
```

Chapter 7

7.10.1 Short Answer

1. (a) 6Ah (b) EAh (c) FDh (d) A9h
2. (a) 9Ah (b) 6Ah (c) 0A9h (d) 3Ah
3. DX = 0002h, AX = 2200h
4. AX = 0306h
5. EDX = 0, EAX = 00012340h
6. The DIV will cause a divide overflow, so the values of AX and DX cannot be determined.
7. DX = 0016h
8. The DIV will cause a divide overflow.
9. In correcting this example, it is easiest to reduce the number of instructions. You can use a single register (ESI) to index into all three variables. ESI should be set to zero before the loop because the integers are stored in little endian order with their low-order bytes occurring first:

```
    mov ecx,8           ; loop counter
    mov esi,0           ; use the same index register
    clc                 ; clear Carry flag
top:
    mov al,byte ptr val1[esi] ; get first number
    sbb al,byte ptr val2[esi] ; subtract second
    mov byte ptr result[esi],al ; store the result
    inc esi             ; move to next pair
    loop top
```

Of course, you could easily reduce the number of loop iterations by adding doublewords rather than bytes.

10. (Shift each bit two positions to the left) = 4080C10140000h

Shown in binary:

```
0001 0000 0010 0000 0011 0000 0100 0000 0101 0000 0000 0000 0000 (before)
0100 0000 1000 0000 1100 0001 0000 0001 0100 0000 0000 0000 0000 (after)
```

7.10.2 Algorithm Workbench

1. Code example:

```
shl  eax,16
sar  eax,16
```

2. Code example:

```
shr  al,1           ; shift AL into Carry flag
jnc  next           ; Carry flag set?
or   al,80h         ; yes: set highest bit
next:                ; no: do nothing
```

3. shl eax,4

4. `shr ebx,2`
5. `ror dl,4` (or: `rol dl,4`)
6. `shld dx,ax,1`
7. This problem requires us to start with the high-order byte and work our way down to the lowest byte:

```
byteArray BYTE 81h,20h,33h
.code
shr byteArray+2,1
rcr byteArray+1,1
rcr byteArray,1
```

8. This problem requires us to start with the low-order word and work our way up to the highest word:

```
wordArray WORD 810Dh,0C064h,93ABh
.code
shl wordArray,1
rcl wordArray+2,1
rcl wordArray+4,1
```

9. Code example:

```
mov ax,3
mov bx,-5
imul bx
mov val1,ax ; product
// alternative solution:
mov al,3
mov bl,-5
imul bl
mov val1,ax ; product
```

10. Code example:

```
mov ax,-276
cwd ; sign-extend AX into DX
mov bx,10
idiv bx
mov val1,ax ; quotient
```

11. Implement the unsigned expression: $val1 = (val2 * val3) / (val4 - 3)$.

```
mov eax,val2
mul val3
mov ebx,val4
sub ebx,3
div ebx
mov val1,eax
```

(You can substitute any 32-bit general-purpose register for EBX in this example.)

12. Implement the signed expression: $val1 = (val2 / val3) * (val1 + val2)$.

```
mov eax,val2
cdq ; extend EAX into EDX
idiv val3 ; EAX = quotient
mov ebx,val1
add ebx,val2
imul ebx
mov val1,eax ; lower 32 bits of product
```

(You can substitute any 32-bit general-purpose register for EBX in this example.)

13. Code example (displays binary value in AX):

```
out16 proc
    aam
    or    ax,3030h
    push eax
    mov  al,ah
    call WriteChar
    pop  eax
    call WriteChar
    ret
out16 endp
```

14. After AAA, AX would equal 0108h. Intel says: First, if the lower digit of AL is greater than 9 or the AuxCarry flag is set, add 6 to AL and add 1 to AH. Then in all cases, AND AL with 0Fh. Here is their pseudocode:

```
IF ((AL AND 0FH) > 9) OR (AuxCarry = 1) THEN
    add 6 to AL
    add 1 to AH
END IF

AND AL with 0FH
```

15. Calculate $x = n \bmod y$, given n and y , where y is a power of 2:

```
.data
dividend DWORD 1000
divisor  DWORD 32      ; must be a power of 2
answer   DWORD ?
.code
mov  edx,divisor      ; create a bit mask
sub  edx,1
mov  eax,dividend
and  eax,edx         ; clear high bits, low bits contain mod value
mov  answer,eax
```

16. Calculate absolute value of EAX without using a conditional jump:

```
mov  edx,eax      ; create a bit mask
sar  edx,31
add  eax,edx
xor  eax,edx
```

17. Shift low bit of AX into high bit of BX:

```
shr  ax,1        ; shift AX into Carry flag
rcr  bx,1        ; shift Carry flag into BX

; Using SHRD:

shrd bx,ax,1
```

18. Count the bits in EBX:

```
CountBits proc
    ; input EBX = integer to evaluate
    ; output EAX = count
    push ebx
    push ecx
    mov  ecx,32      ; loop counter
    mov  eax,0      ; counts the '1' bits
L1: shr  ebx,1      ; shift into Carry flag
```

```

    jnc     L2           ; Carry flag set?
    inc     eax         ; yes: add to bit count
L2: loop   L1           ; continue loop until done
    pop     ecx
    pop     ebx
    ret
Countbits endp

```

19. Solution code:

```

    mov     ecx,32      ; loop counter
    mov     bl,0        ; counts the '1' bits
L1: shr    eax,1       ; shift into Carry flag
    jnc     L2         ; Carry flag set?
    inc     bl         ; yes: add to bit count
L2: loop   L1         ; continue loop

; if BL is odd, clear the parity flag
; if BL is even, set the parity flag

    shr    bl,1
    jc     odd
    mov    bh,0
    or     bh,0        ; PF = 1
    jmp   next
odd:
    mov    bh,1
    or     bh,1        ; PF = 0
next:

```

Chapter 8

8.10.1 Short Answer

1. Code example:

```
mov esp,ebp
pop ebp
```

2. EAX

3. It passes an integer constant to the RET instruction. This constant is added to the stack pointer right after the RET instruction has popped the procedure's return address off the stack.

4. LEA can return the offset of an indirect operand; it is particularly useful for obtaining the offset of a stack parameter.

5. Four bytes

6. The C calling convention allows for variable-length parameter lists.

7. False

8. False

9. True (because the immediate value will be interpreted as an address)

8.10.2 Algorithm Workbench

1. Stack frame:

10h	[EBP + 16]
20h	[EBP + 12]
30h	[EBP + 8]
(return addr)	[EBP + 4]
EBP	<--ESP

2. Code example:

```
AddThree PROC
; modeled after the AddTwo procedure
push ebp
mov  ebp,esp
mov  eax,[ebp + 16]      ; 10h
add  eax,[ebp + 12]     ; 20h
add  eax,[ebp + 8]      ; 30h
pop  ebp
ret  12
AddThree ENDP
```

3. Declaration: LOCAL pArray:PTR DWORD

4. Declaration: LOCAL buffer[20]:BYTE

5. Declaration: LOCAL pwArray:PTR WORD

6. Declaration: LOCAL myByte:SBYTE
7. Declaration: LOCAL myArray[20]:DWORD
8. Code example:

```
SetColor PROC USES eax,  
    forecolor:BYTE, backcolor:BYTE  
  
    movzx eax,backcolor  
    shl  eax,4  
    add  al,forecolor  
    call SetTextColor  
  
    ret  
SetColor ENDP
```

9. Code example:

```
WriteColorChar PROC USES eax,  
    char:BYTE,forecolor:BYTE, backcolor:BYTE,  
  
    INVOKE SetColor, forecolor, backcolor  
    mov  al,char  
    call WriteChar  
    ret  
WriteColorChar ENDP
```

10. Code example:

```
DumpMemory PROC USES esi ebx ecx,  
    address:DWORD,           ; starting address  
    units:DWORD,           ; number of units  
    unitType:DWORD         ; unit size  
  
    mov  esi,address  
    mov  ecx,units  
    mov  ebx,unitType  
    call DumpMem  
    ret  
DumpMemory ENDP
```

11. Code example:

```
MultArray PROC USES esi ebx ecx,  
    array1:PTR DWORD, array2:PTR DWORD,  
    count:DWORD  
  
MultArray PROTO,  
    array1:PTR DWORD, array2:PTR DWORD,  
    count:DWORD
```

Chapter 9

9.9.1 Short Answer

1. 1 (set)
2. 2 is added to the index register
3. Regardless of which operands are used, CMPS still compares the contents of memory pointed to by ESI to the memory pointed to by EDI.
4. 1 byte beyond the matching character.
5. REPNE (REPZ).
6. 1 (set)
7. JNE is used to exit the loop and insert a null byte into the string when no more characters are to be trimmed.
8. The digit is unchanged.
9. REPNE (REPZ).
10. The length would be $(EDI_{\text{final}} - EDI_{\text{initial}}) - 1$.
11. The maximum comparisons for 1,024 elements is 11.
12. The Direction flag is cleared so that the STOSD instruction will automatically increment the EDI register. Instead, if the flag were set, EDI would decrement and move backwards through the array.
13. EDX and EDI were already compared.
14. Change each JMP L4 instruction to JMP L1.

9.9.2 Algorithm Workbench

1. [ebx + esi]
2. array[ebx + esi]
3. Code example:

```
mov esi,2           ; row
mov edi,3           ; column
mov eax,[esi*16 + edi*4]
```

4. CMPSW example:

```
mov ecx,count
mov esi,offset sourcew
mov edi,offset targetw
cld
repe cmpsw
```

5. SCASW example:

```
cld
mov   ecx,4
mov   edi,offset wordArray
repne scasw
sub   edi,4           ; adjust the offset
mov   eax,edi
```

6. Str_compare example:

```
.data
str1 byte "ABCDE",0           ; larger string
str2 byte "ABCCD",0
.code
    invoke Str_compare, offset str1, offset str2
    ja   strlbigger
    mov  edx,offset str2
    jmp  print
strlbigger:
    mov  edx,offset str1
print:
    call WriteString
    call Crlf
```

7. Str_trim example:

```
.data
prob7string byte "ABCD@@@",0
.code
    invoke Str_trim, offset prob7string, '@'
```

8. Str_lcase example (converts string to lower case):

```
Str_lcase PROC USES eax esi,
    pString:PTR BYTE
    mov esi,pString
L1:
    mov al,[esi]           ; get char
    cmp al,0               ; end of string?
    je  L3                 ; yes: quit
    cmp al,'A'             ; below "A"?
    jb  L2                 ; below "A"?
    cmp al,'Z'             ; above "Z"?
    ja  L2                 ; above "Z"?
    or  BYTE PTR [esi],00100000b ; convert to lower case
L2:inc esi ; next char
    jmp L1
L3: ret
Str_lcase ENDP
```

9. 64-bit Str_trim procedure:

```
-----
Str_trim PROC uses rax rcx rdi
; Removes all occurrences of a given character from
; the end of a string.
; Receives: RCX points to the string, AL contains the trim character
; Returns: nothing
-----
.data
strtrimchar byte ?
.code
    mov strtrimchar,al
    mov rdi,rcx           ; save pointer to string
    call Str_length       ; puts length in RAX
```

```

    cmp rax,0           ; length zero?
    je  L3             ; yes: exit now
    mov rcx,rax        ; no: RCX = string length
    dec rax
    add rdi,rax        ; point to null byte at end

L1: mov  al,[rdi]      ; get a character
    cmp  al,strtrimchar ; character to be trimmed?
    jne  L2            ; no: insert null byte
    dec  rdi           ; yes: keep backing up
    loop L1            ; until beginning reached

L2: mov  BYTE PTR [rdi+1],0 ; insert a null byte
L3:  ret
Str_trim ENDP

```

10. Base-index operand in 64-bit mode:

```
array[rsi*TYPE array]
```

11. Two-dimensional array indexing, 32-bit mode:

```
mov eax,table[ebx + edi*TYPE myArray]
```

12. Two-dimensional array indexing, 64-bit mode:

```
mov rax,table[rbx + rdi*TYPE myArray]
```

Chapter 10

10.7.1 Short Answer

1. Structures are essential whenever you need to pass a large amount of data between procedures. One variable can be used to hold all the data.
2. Answers:
 - a. Yes.
 - b. No.
 - c. Yes.
 - d. Yes.
 - e. No.
3. False.
4. To permit the use of labels in a macro that is invoked more than once by the same program.
5. ECHO (also, the %OUT operator, which is shown later in the chapter).
6. ENDEF.
7. List of relational operators:
 - LT Less than
 - GT Greater than
 - EQ Equal to
 - NE Not equal to
 - LE Less than or equal to
 - GE Greater than or equal to
8. The substitution (&) operator resolves ambiguous references to parameter names within a macro.
9. The literal-character operator (!) forces the preprocessor to treat a predefined operator as an ordinary character.
10. The expansion operator (%) expands text macros or converts constant expressions into their text representations.

10.7.2 Algorithm Workbench

1. Structure definition:

```
MyStruct STRUCT
    field1 WORD ?
    field2 DWORD 20 DUP(?)
MyStruct ENDS
```


2. Code example:

```
.data
time SYSTEMTIME <>
.code
mov ax,time.wHour
```

3. Code example:

```
myShape Triangle < <0,0>, <5,0>, <7,6> >
```

4. Declare and initialize an array of Triangle structures:

```
.data
ARRAY_SIZE = 5
triangles Triangle ARRAY_SIZE DUP(<>)
.code
    mov ecx,ARRAY_SIZE
    mov esi,0
L1: mov eax,11
    call RandomRange
    mov triangles[esi].Vertex1.X, ax
    mov eax,11
    call RandomRange
    mov triangles[esi].Vertex1.Y, ax
    add esi,TYPE Triangle
    loop L1
```

5. Code example:

```
mPrintChar MACRO char,count
LOCAL temp

.data
temp BYTE count DUP(&char),0

.code
push edx
mov edx,OFFSET temp
call WriteString
pop edx
ENDM
```

6. Code example:

```
mGenRandom MACRO n
mov eax,n
call RandomRange
ENDM
```

7. mPromptInteger:

```
mPromptInteger MACRO prompt,returnVal
mWriteprompt
call ReadInt
mov returnVal,eax
ENDM
```

8. Code example:

```
mWriteAt MACRO X,Y,literal
mGotoxy X,Y
mWrite literal
ENDM
```

9. Code example:

```
mWriteStr namePrompt
1 push edx
1 mov  edx,OFFSET namePrompt
1 call WriteString
1 pop  edx
```

10. Code example:

```
mReadStr customerName
1 push ecx
1 push edx
1 mov  edx,OFFSET customerName
1 mov  ecx,(SIZEOF customerName) - 1
1 call ReadString
1 pop  edx
1 pop  ecx
```

11. Code example:

```
;-----
mDumpMemx MACRO varName
;
; Displays a variable in hexadecimal, using the
; variable's attributes to determine the number
; of units and unit size.
;-----
push ebx
push ecx
push esi
mov esi,OFFSET varName
mov ecx,LENGTHOF varName
mov ebx,TYPE varName
call DumpMem
pop esi
pop ecx
pop ebx
ENDM

; Sample calls:
.data
array1 BYTE 10h,20h,30h,40h,50h
array2 WORD 10h,20h,30h,40h,50h
array3 DWORD 10h,20h,30h,40h,50h
.code
mDumpMemx array1
mDumpMemx array2
mDumpMemx array3
```

12. Macro using a default argument initializer:

```
mWriteLn MACRO text:=<" ">
mWrite text
call Crlf
ENDM
```

13. Macro using IF, ELSE, ENDIF:

```
mCopyWord MACRO intVal
IF (TYPE intVal) EQ 2
mov ax,intVal
ELSE
ECHO Invalid operand size
ENDIF
ENDM
```

14. Macro using IF to check value of a parameter

```
mCheck MACRO Z
IF Z LT 0
ECHO **** Operand Z is invalid ****
ENDIF
ENDM
```

15. Macro uses the & operator when parameter is embedded in a string:

```
CreateString MACRO strVal
.data
temp BYTE "Var&strVal",0
.code
ENDM
```

16. Source code generated by the mLocate macro:

```
mLocate -2,20
;(no code generated because xval < 0)
mLocate 10,20
1 mov bx,0
1 mov ah,2
1 mov dh,20
1 mov dl,10
1 int 10h
mLocate col,row
1 mov bx,0
1 mov ah,2
1 mov dh,row
1 mov dl,col
1 int 10h
```

Chapter 11

11.7.1 Short Answer

1. BOOL = byte, COLORREF = DWORD, HANDLE = DWORD, LPSTR = PTR BYTE, WPARAM = DWORD.
2. GetStdHandle.
3. ReadConsole.
4. The COORD structure contains X and Y screen coordinates in character measurements.
5. SetFilePointer.
6. SetConsoleTitle.
7. SetConsoleScreenBufferSize.
8. SetConsoleCursorInfo.
9. SetConsoleTextAttribute.
10. WriteConsoleOutputAttribute.
11. Sleep.
12. (A program that calls CreateWindowEx is shown in Section 11.2.6.) The prototype for CreateWindowEx is located in the GraphWin.inc file:

```
CreateWindowEx PROTO,  
classexWinStyle:DWORD,  
className:PTR BYTE,  
winName:PTR BYTE,  
winStyle:DWORD,  
X:DWORD,  
Y:DWORD,  
rWidth:DWORD,  
rHeight:DWORD,  
hWndParent:DWORD,  
hMenu:DWORD,  
hInstance:DWORD,  
lpParam:DWORD
```

The fourth parameter, winStyle, determines the window's style characteristics. In the WinApp.asm program in Section 11.2.6, when we call CreateWindowEx, we pass it a combination of predefined style constants:

```
MAIN_WINDOW_STYLE = WS_VISIBLE + WS_DLGFRAME + WS_CAPTION  
+ WS_BORDER + WS_SYSMENU + WS_MAXIMIZEBOX + WS_MINIMIZEBOX  
+ WS_THICKFRAME
```

The window described here will be visible, and it will have a dialog box frame, a caption bar, a border, a system menu, a maximize icon, a minimize icon, and a thick surrounding frame.

13. Choose any two of the following (from GraphWin.inc):

```
MB_OK, MB_OKCANCEL, MB_ABORTRETRYIGNORE, MB_YESNOCANCEL, MB_YESNO,  
MB_RETRYCANCEL, MB_CANCELTRYCONTINUE
```

14. Icon constants (choose any two):

```
MB_ICONHAND, MB_ICONQUESTION, MB_ICONEXCLAMATION, MB_ICONASTERISK
```

15. Tasks performed by WinMain (choose any three):

- Get a handle to the current program.
- Load the program's icon and mouse cursor.
- Register the program's main window class and identify the procedure that will process event messages for the window.
- Create the main window.
- Show and update the main window.
- Begin a loop that receives and dispatches messages.

16. The WinProc procedure receives and processes all event messages relating to a window. It decodes each message, and if the message is recognized, carries out application-oriented (or application-specific) tasks relating to the message.

17. The following messages are processed:

- WM_LBUTTONDOWN, generated when the user presses the left mouse button.
- WM_CREATE, indicates that the main window was just created.
- WM_CLOSE, indicates that the application's main window is about to close.

18. The ErrorHandler procedure, which is optional, is called if the system reports an error during the registration and creation of the program's main window.

19. The message box is shown before the application's main window appears.

20. The message box appears before the main window closes.

21. A linear address is a 32-bit integer ranging between 0 and FFFFFFFh, which refers to a memory location. The linear address may also be the physical address of the target data if a feature called paging is disabled.

22. When paging is enabled, the processor translates each 32-bit linear address into a 32-bit physical address. A linear address is divided into three fields: a pointer to a page directory entry, a pointer to a page table entry, and an offset into a page frame.

23. The linear address is automatically a 32-bit physical memory address.

24. Paging makes it possible for a computer to run a combination of programs that would not otherwise fit into memory. The processor does this by initially loading only part of a program in memory while keeping the remaining parts on disk.

25. The LDTR register.

26. The GDTR register.

27. One.

28. Many (each task or program has its own local descriptor table).

29. Choose any four from the following list: base address, privilege level, segment type, segment present flag, granularity flag, segment limit.

30. Page Directory, Page Table, and Page (page frame).

31. The Table field of a linear address (see Figure 11-4).
32. The Offset field of a linear address (see Figure 11-4).

11.7.1 Algorithm Workbench

1. Example from the ReadConsole.asm program in Section 11.1.4:

```
INVOKE ReadConsole, stdInHandle, ADDR buffer,
BufSize - 2, ADDR bytesRead, 0
```

2. Example from the Console1.asm program in Section 11.1.5:

```
INVOKE WriteConsole,
consoleHandle,          ; console output handle
ADDR message,          ; string pointer
messageSize,          ; string length
ADDR bytesWritten,    ; returns num bytes written
0                      ; not used
```

3. Calling CreateFile when reading an input file:

```
INVOKE CreateFile,
ADDR filename,          ; ptr to filename
GENERIC_READ,          ; access mode
DO_NOT_SHARE,          ; share mode
NULL,                  ; ptr to security attributes
OPEN_EXISTING,         ; file creation options
FILE_ATTRIBUTE_NORMAL, ; file attributes
0                      ; handle to template file
```

4. Calling CreateFile to create a new file:

```
INVOKE CreateFile,
ADDR filename,
GENERIC_WRITE,
DO_NOT_SHARE,
NULL,
CREATE_ALWAYS,
FILE_ATTRIBUTE_NORMAL,
0
```

5. Calling ReadFile:

```
INVOKE ReadFile,          ; read file into buffer
fileHandle,
ADDR buffer,
bufSize,
ADDR byteCount,
0
```

6. Calling WriteFile:

```
INVOKE WriteFile,        ; write text to file
fileHandle,              ; file handle
ADDR buffer,             ; buffer pointer
bufSize,                 ; number of bytes to write
ADDR bytesWritten,       ; number of bytes written
0                        ; overlapped execution flag
```

7. Calling MessageBox:

```
INVOKE MessageBox, hMainWnd, ADDR GreetText,  
ADDR GreetTitle, MB_OK
```

Chapter 12

12.6.1 Short Answer

1. $1101.01101 = 13/1 + 1/4 + 1/8 + 1/32$
2. 0.2 generates an infinitely repeating bit pattern.
3. $11011.01011 = 1.101101011 \times 2^4$
4. $0000100111101.1 = 1.001111011 \times 2^{28}$
5. Quiet NaN and Signaling NaN
6. REAL10 80 bits
7. It pops ST(0) off the stack.
8. FCHS.
9. None, m32fp, m64fp, stack register
10. FISUB converts the source operand from integer to floating-point.
11. FCOM, or FCOMP
12. Code example:

```
fnstsw ax
lahf
```
13. FILD
14. RC field

12.6.2 Algorithm Workbench

1. $+1110.011 = 1.110011 \times 2^3$, so the encoding is 0 1000010 1100110000000000000000
2. $5/8 = 0.101$ binary
3. $17/32 = 0.10001$ binary
4. $+10.75 = +1010.11 = +1.01011 \times 2^3$, encoded as 0 1000010 0101100000000000000000
5. $-76.0625 = -01001100.0001 = -1.0011000001 \times 2^6$, encoded as:
1 10000101 001100000100000000000000
6. Code example:

```
fnstsw ax
lahf
```
7. 1.010101101
8. 1.010101101 rounded to nearest even becomes 1.010101110.

9. Assembly language code:

```
.data
B REAL8 7.8
M REAL8 3.6
N REAL8 7.1
P REAL8 ?
.code
fld M
fchs
fld N
fadd B
fmul
fst P
```

10. Assembly language code:

```
.data
B DWORD 7
N REAL8 7.1
P REAL8 ?
.code
fld N
fsqrt
fiadd B
fst P
```

11. (a) 8E (b) 8A (c) 8A (d) 8B (e) A0 (f) 8B

12. (a) 06 (b) 56 (c) 1D (d) 55 (e) 84 (f) 81

13. Machine language bytes:

```
a. 8E D8
b. A0 00 00
c. 8B 0E 01 00
d. BA 00 00
e. B2 02
f. BB 00 10
```

Chapter 13

13.7 Review Questions

1. The memory model determines whether near or far calls are made. A near call pushes only the 16-bit offset of the return address on the stack. A far call pushes a 32-bit segment/offset address on the stack.
2. C and C++ are case sensitive, so they will only execute calls to procedures that are named in the same fashion.
3. Yes, many languages specify that EBP (BP), ESI (SI), and EDI (DI) must be preserved across procedure calls.
4. Yes.
5. No.
6. No.
7. A program bug might result because the `__fastcall` convention allows the compiler to use general-purpose registers as temporary variables.
8. Use the LEA instruction.
9. The LENGTH operator returns the number of elements in the array specified by the DUP operator. For example, the value placed in EAX by the LENGTH operator is 20:

```
myArray DWORD 20 DUP(?), 10, 20, 30
.code
mov eax,LENGTH myArray ; 20
```
10. The SIZE operator returns the product of TYPE (4) * LENGTH.
11. `printf` `PROTO C, pString:PTR BYTE, args:VARARG.`
12. X will be pushed last.
13. To prevent the decoration (altering) of external procedure names by the C++ compiler. Name decoration (also called name mangling) is done by programming languages that permit function overloading, which permits multiple functions to have the same name.
14. If name decoration is in effect, an external function name generated by the C++ compiler will not be the same as the name of the called procedure written in assembly language. Understandably, the assembler does not have any knowledge of the name decoration rules used by C++ compilers.
15. Virtually no changes at all, showing that array subscripts can be just as efficient as pointers when manipulating arrays.